Trinity College Dublin Advanced Topics in Quantum Information Theory Hilary Term 2024 Quantum Information Theory Group

# Exercises - Week 27

## Problem 1: Bell states and maximal entanglement

The four Bell states form a basis for the Hilbert space of two qubits. In this exercise we will further explore why they are called 'maximally entangled'.

- (i) Write down the four Bell states and consider local unitary operations of the form  $U_{AB} = U_A \otimes U_B$ . Can any Bell state converted into any other by such a local unitary operation? If no, provide a proof. If yes, specify a sufficient set of local unitary operations.
- (ii) Can a product state  $|\psi\rangle_{AB} = |a\rangle \otimes |b\rangle$  be obtained from a Bell state by a local unitary operation?
- (iii) It is important to know that quantum teleportation does not permit faster than light communication. Proof this by showing that in the teleportation protocol introduced in the lecture, Bob does not obtain any information about Alice' input state  $|\psi\rangle$  until she has communicated her measurement result. Why is your result is not in conflict with the functioning of quantum teleportation?
- (iv) Show that Alice can use an extension of the teleportation protocol from the lecture to create any shared state  $|\sigma\rangle_{AB}$  between her laboratory and Bob's by using one shared Bell state  $|\phi^+\rangle$ . *Hint:* Consider Alice and Bob to start with the initial state  $|\sigma\rangle_{12} \otimes |\phi^+\rangle_{AB}$  where Alice

controls laboratories 1, 2, and A; Bob controls laboratory B. As in the original teleportation protocol encountered in the lecture, Alice and Bob may only perform measurements and unitary operations on systems in their control, as well as classical communication.

(v) We have called the Bell states 'maximally entangled'. How is this terminology justified for pure two-qubit states in light of what we have found out in this exercise?

### Problem 2: QRTs from basic postulates

In the lecture, we have stated the *free operations postulate (FOP)* as a postulate. In this exercise, we will show that it can in fact be derived from other basic postulates for quantum resource theories (QRTs), namely

- (free preparation of free states) The set of free states  $\mathcal{F}$  is in one-to-one correspondence with the set of free preparation maps  $\Omega_{\phi} : \mathbb{C} \to \mathcal{H} \ (\forall \mathcal{H})$ . In other words,  $\phi$  is a free state if and only if  $\mathcal{O}$  includes a free map  $\Omega_{\phi} : \mathbb{C} \to \mathcal{H}$  with  $\Omega(1) = \phi$ .
- (free identity operations) For any permitted Hilbert space  $\mathcal{H}$ , the associated identity operation is a free operation:  $id_{\mathcal{H}} \in \mathcal{O}$ ,
- (free concatenation of free operations)  $\Omega, \Omega' \in \mathcal{O} \Rightarrow \Omega' \circ \Omega \in \mathcal{O}$ , for all  $\Omega$  and  $\Omega'$  for which the output Hilbert space of  $\Omega$  matches the input Hilbert space of  $\Omega'$  and where  $\circ$  stands for successive application.
- (i) Write down the FOP. (ii) Then, derive it from the above.

### Problem 3: Separable states and free operations

For an *N*-partite Hilbert space, recall the set of separable states  $\mathcal{F}_s$  and the set of separable operations  $(\mathcal{O}_{sep})$ .

- (i) Consider O<sub>c-max</sub> induced by F<sub>s</sub>. Show that O<sub>sep</sub> ⊆ O<sub>c-max</sub> i.e., there is no operation in O<sub>sep</sub> which is not also in O<sub>c-max</sub>.
  *Remark*: In fact, O<sub>sep</sub> = O<sub>c-max</sub> but you don't need to prove the other direction i.e., that there is no operation in O<sub>c-max</sub> which is not also in O<sub>sep</sub>.
- (ii) Consider the case of bipartite separable states  $\mathcal{F}_{2s}$ . Show that any bipartite separable state can be written as a classical mixture of pure, orthonormal, bipartite product states (i.e., states of the form  $|a\rangle_A |b\rangle_B$ ). Is any classical mixture of product states a separable state? (*Remark:* since this technically includes mixtures with a single component, note that the only separable pure states are product states).
- (iii) Consider again  $\mathcal{F}_{2s}$ . The set  $\mathcal{O}_{max}$  induced by  $\mathcal{F}_{2s}$  is the set of non-entangling operations  $\mathcal{O}_{ne}$ . The swap operation SWAP is implicitly defined by  $SWAP(|a\rangle_A |b\rangle_B) = |b\rangle_A |a\rangle_B$  for any pair of systems and states. Recalling that  $\mathcal{O}_{c-max} = \mathcal{O}_{sep}$ . Show that for  $\mathcal{F}_{2s}$ ,  $\mathcal{O}_{c-max} \subsetneq \mathcal{O}_{max}$  by proofing that  $SWAP \in \mathcal{O}_{ne}$  and  $SWAP \notin \mathcal{O}_{sep}$ .

### **Problem 4: Properties of entanglement monotones**

In this exercise we will examine some of properties and desiderata for entanglement monotones more closely by focusing on specific examples.

- (i) Show that the Schmidt rank is an entanglement monotone for pure states. How can we turn the Schmidt rank from a monotone into a measure (i.e., by making it weakly discriminant)? Is this derived measure faithful?
- (ii) We can formally construct a **resource monotone from a distance function** d by defining:

$$E_d(\rho) := \inf_{\varphi \in \mathcal{F}} d(\rho, \varphi)$$

Note that any  $E_d$  defined in this way is faithful.

Show that if d obeys a data processing inequality  $d(\rho, \sigma) \ge d(\Lambda(\rho), \Lambda(\sigma))$  for all quantum channels  $\Lambda$  (not just free ones) then  $E_d$  is monotonic under all non-entangling operations. *Remark:* Recall the requirements that a distance function has to fulfill:

- $d(x,y) \ge 0$  (non-negativity)
- $d(x,y) = 0 \Leftrightarrow x = y$  (identity of indiscernibles)
- d(x,y) = d(y,x) (symmetry)
- $d(x,z) \le d(x,y) + d(y,z)$  (triangle inequality)
- (iii) Recall the **relative entropy**  $R(\rho||\sigma) := -S(\rho) \text{Tr}[\rho \log \sigma]$ . R is not a distance on state space because it is neither symmetric  $R(\rho||\sigma) \neq R(\sigma||\rho)$  nor does it satisfy the triangle inequality. This is equally true of its classical variant (called Kullback-Leibler divergence).

a) Show that R is not symmetric by finding two states,  $\rho$  and  $\sigma$  as an example.

Regardless of this, R can still be used as an entanglement monotone  $E_R$  using the same construction as for  $E_d$  as above.

b) The quantum relative entropy obeys joint convexity in its arguments, i.e., for p+q=1and states  $\rho, \sigma, \alpha, \beta$ ::

 $R(p\rho + q\sigma || p\alpha + q\beta) \le pR(\rho || \alpha) + qR(\sigma || \beta)$ 

Use this to show that  $E_R$  is convex.