Recap: • Linear map $\mathcal{E}: L(\mathcal{H}_1) \rightarrow L(\mathcal{H}_2) \iff \operatorname{wahrix} \mathcal{E}_{\mathcal{L}} \in L(\mathcal{H}_1, \otimes \mathcal{H}_2)$ E = (E @ I,) [H>< M], where M> E H = M is the maxinally entangled state. $\mathcal{E}: L(\mathcal{H}_1) \rightarrow L(\mathcal{H}_1)$ is CPTP $\iff \mathcal{E}_c \gg \mathcal{O} \quad \mathcal{E} \quad \mathcal{H}_{2'}[\mathcal{E}_c] = \mathbf{1}_{\mathcal{L}}$ "=>" By direct insertion " <= " By proving that E > 0 leads to a Krans de composition. Note: E is CP => E >0 => Kraus decomposition 2, E is CP if and only if E admits a knows decomposition above proof. ("everything is Choi mabrices") Examples :

o, partially dephasing map $\mathcal{E}: \mathcal{L}(\mathcal{H}_{1}) \rightarrow \mathcal{L}(\mathcal{H}_{1}) \cong \mathcal{L}(\mathcal{H}_{1})$ $\mathcal{E}[S] = p \ tr[S] \stackrel{\text{def}}{=} tr (1-p)S$ $L_{2} \mathcal{E}_{c} = p \frac{\mathcal{I}_{1}}{\mathcal{A}} \mathfrak{G} \mathcal{I}_{1} + (1-p) [\mathcal{H}_{2} \mathcal{H}]$ í S Choi stete of tr channel change b, let R: C -> L(7t.) be a map that satisfies R[j] > 0 4320 (complete positivity) show this (trace preservation) and tr [R[]] = 3 The map R is called a preparation with comes. pondig Cho: state R_=: SEL(H, BE) = L(H,) it sakisfies 320 and try [3] = hr[s] = 1 La Quantum States are Choi matrices of preperation maps.

c, Let F: L(H_) → L(H_1) = C be a line~ map from mahiers to munders flick satisfies J[y] = 0 ty = 0 (complete posihisty) show this and tr[f[y]] = tr[y] (trace non-increasing) F is called an effect. Its choi state Fe L(H, OC) = L(H,) salisfies $F \ge 0$, $t_c[F] = F \le 4_1$ if \overline{f} is trace preserving, then $\overline{F} \ge 0$, $\overline{fr} \subseteq \overline{f} = 4$ 4.4 Choi representation and the non-vinqueness of the Krans decomposition Note: The Chai makin E of a mop E is unique but its Krans decomposition is uor. La How many Urans oprotors are required (ct

most), and how are different trans represubabions velocked?

Every CP map E: L(74,) -> L(74,) Thm. requires at most dat da Krans operators {K, } and they can be chose such flat $tr(k_{\alpha}^{\dagger}k_{\beta}) = S_{\alpha\beta}tr(k_{\alpha}^{\dagger}k_{\beta})$ tx, > conduical Krans deco-positio-

Proof: (for d, = d,'=: d). Let E: L(H,) -> L(H') and E = E], (2>2x). Toke the Krows opentors Ky = E Az cildscil from the previous theorem. There is at most d² of them and it is easy to check that they satisfy $tr(K_{2}^{\dagger}K_{p}) = S_{ap} tr(K_{a}^{\dagger}K_{p})$ U

Def.: (Rank of a CP map): The rank of a CP map 's the mininal much of Kraus operators it requires for its representation.

Lemma: The rank of a CP map & coincides with the rank of its Choi matrix E.

Proof: From above, we know I know decomposition Ek, 3 with D = rank (Ec). Now, let {kip} be another Kraus decomposition of E. Then; $\mathcal{E}_{c} = \sum_{\beta=1}^{\infty} |\mathbf{k}_{\beta}'\rangle < |\mathbf{k}_{\beta}'|$, with $|\mathbf{k}_{\beta}'\rangle = (\mathbf{k}_{\beta} \otimes \mathbf{1})|\mathbf{H}\rangle$.

=> Since rank (E.) 'is the minimal number of pure states required for such a decomposition of Ec, we have rank (Ec) & D'. 21 rank (Ec) is the veiniel number of required Kraus operators.

La Important: Every choice of krows operators corresponds to a squarerook X of E = X X S All Krans operators are of the form Kx = Z < i | Kx >< i | , where E = J | Kx >< Kx |.

L> The canonical representation stemming from the mique positive squarerook X = 5 JZ, 1x><x1 of E is called mininal.

(Kraus de compositions are related by isometries") Lemma: Denote the mininal krows decomposition of E by EK, 3. All other Krows decompositions EKp? of satisfy $K'_{\beta} = \sum_{\alpha} V_{\beta\alpha} K_{\alpha},$ where V is an isometry, i.e. Vt V = 11.

La Characteization of all Kraus decompositions. Proof: Let X be the migne positive squareroot of Ec and let t' be some other squareroot of Ec. All squareroots of E are related to X by isometry, i.e. $\exists V$ with $V^{\dagger}V = A$ and X' = VX (show this).

Now, let {Ky} be the Krans operators "stamming from" X and let { kp} be those skenning from X'. $=> |k_{\mu}'\rangle_{\mathcal{F}} = \chi_{\mu\mathcal{F}}' = (V\chi)_{\mu\mathcal{F}}' = \sum_{\alpha} V_{\mu\alpha} |k_{\alpha}\rangle_{\mathcal{F}}$

 $=> |k_{\beta}'> = \sum_{\alpha} V_{\beta\alpha} |k_{\alpha}>$ $L_{>}$ $K_{p}' = \sum_{\alpha} V_{p\alpha} K_{\alpha}$

